

HiPAS High Performance Adaptive Schema Migration with Minimum Downtime Option

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Background



- Minimum Downtime Schema Migration and Continous Replication
 - needed very often
 - business and data critical
 - high demand of intensive planning
- Implemented completely in PL /SQL
 - adding up the best practices from Data Pump, O2O, Golden Gate
 - only one PL/SQL package on source and destination
- Academic approach
 - Self Adaptive (artificial intelligence)
 - Developed together with the University of Potsdam/Berlin

Agenda





Migration Challenges



- Short Downtime
 - expensive unavailability due to opportunity costs
- Storage I/O Controller Utilization
 - average utilization of 70% as optimal [2]
 - table diversity (empty, small, very large), up to 70,000 tables
- Endianness
 - byte order changes, e.g., from Solaris to Linux





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Migration Approach Differentiation



- Invocation layer
 - Storage
 - OS
 - Database
- change of platform
- change of endianness
- change of character set
- Downtime proportionality
 - Size of migration data
 - Data alteration rate

Migration Method	Invocation Layer/ Granularity	Downtime Proportionality	Platform Change	Endianness Change	Character Set Change
Storage Replication	Storage/ Storage	negligible	no	no	no
Transportable Database	OS/ Database	Database Size	yes	no	no
Transportable Tablespaces	OS/ Tablespace	Tablespace Size	yes	no	no
Cross Platform Transportable Tablespace	OS/ Tablespace	Tablespace Size	yes	yes	no
Transportable Tablespaces using Cross Platform Incremental Backups	OS/ Tablespace	Data Alteration Rate	yes	no	no
Oracle-to-Oracle (O2O)	OS/ Schema	Amount of Migration Data	yes	yes	no
Datapump	Database/ Value	Amount of Migration Data	yes	yes	yes
Export/Import	Database/ Value	Amount of Migration Data	yes	yes	yes

Prior Analysis



Average Structure of Allocated Data(based on 41 productively running SAP Systems)



 irrelevant data can be excluded when migrating on logical database layer

HiPAS Architecture



Everything is a tuple



Adaptive Data Transfer



- Enabling adaptive behavior during transfer phase
 - partitioning into equally sized transfer bundles
 - Number of running transfer jobs can be reduced or increased
- Two approaches were developed and evaluated
 - Adaption: based on an incremental adjustment process, until changes do not evoke further improvements, thus, reaching the state of an optimal parallelization degree
 - Anticipation: makes continuously new modification decisions independently of each other, based on knowledge about used and monitored resources



Self-Adaptive Software



"Self-adaptive software evaluates its own behavior and **changes behavior when** the evaluation indicates that it is not accomplishing what the software is intended to do, or when **better** functionality or **performance is possible**." [4]

"Self-adaptive software **modifies** its own **behavior in response to** changes in its operating **environment**. [...]" [5]

- Self-Properties of self-adaptive software [6]
 - Self-configuring
 - Self-healing
 - Self-optimizing
 - Self-protecting

Design Space Dimensions



Observation Presentation Environment-Awareness Concurrency events Storage System Average write time CPU Average read time Memory Redo log buffer size Self-Awareness Available memory size Number of running jobs • etc. Control Identification and Enabling Adaption Master/slave control Plugin architecture pattern in distributed Table MIG Control as system interface

[7]

Adaptive Capabilities of HiPAS



"Optimizer" plugin for data transfer phase
acts according to MAPE feedback loop [6]



Database System



Control Dimension



Monitor, Analyse and Plan





Evaluation



- Adaption of parallelization degree according to system environment and migration data
- ~ 123 MByte/s per 1 gbit network interface
- ~ 1 GByte/s per 10 gbit network interface



How does it work?





How does it work?



- PL/SQL only
- SQLNET only
 - no temporary Storage necessary
- Source and Destination RAC aware
 - automatic multi instance parallelization
- Everything protected by Oracle transactional integrity
 - no data loss possible
 - Restart after failure / server / network outage
 - automatically
 - no Re-copy of row sets
- Parallel Index Build

How does it work?



- dbms_metadata on source
- Stats on source
- create table extents on dest
- PL/SQL Long to LOB conversion on source
- University Solution for transfer parallelization
- create dbms_scheduler jobs
- transfer table rows, LOBs
- calibrate IO / Auto DOP for indexing on dest
- Count rows and select "source" minus "dest"
- generate compliance report





- non-adaptive and sequential migrations leave useful resources idle or need to be tuned manually
 "self adaptive is always better"
- Iogical transfer
 - platform, version, endianess and character set independent
- Ultra Fast parallel LOB interface
- Copy Performance of 3 to 5 TByte per Hour
 - adequate Network and I/O Bandwith necessary
- Easy Fallback source stays untouched

Conclusions



- Remap everything
 - User
 - Tablespaces
 - Table / Tablespace Mapping
 - create object attributes
 - Index table compression
- Compliance Check
- Diff Report for rows and metadata

Minimum Downtime Option



- works without EE or Partitioning
- provides same functionality and benefits
 - easy fallback
 - protected by oracle transactional integrity
 - Remap everything
 - Diff Report for all rows and metadata

Minimum Downtime - Capture



- Capture changes while transfer base data
 - List of Transactions
 - ^I Trigger
 - generate list of changes SCN based
 - Old Value / New Value / SCN / ID
 - Uses Log Stream to doublecheck
- Generates List of Sqls
 - Capture / Apply to other DB Platforms possible
- Parallel Capture and Apply

Minimum Downtime – SCN Copy PASOLFORA

Dirty Read Option dismissed

- " "Dirty" Reads (different SCNs per Rowset)
 - merge changes at the end of transfer
 - Row need apply / Row newer than change
 - like Oracle Recovery

Select ... as of (same SCN for all Rowsets)

- Undo Guarantee
- generates insert sqls for multi DB Plattform
- Trigger on Large Tables
- Small Tables in switchover downtime
- apply list of changes ordered

Replication



- Initial Load by Hipas Base Schema Transfer
- Replication based on Hipas capture
- Trigger based
 - thin and fast implementation (rac aware)
 - blacklist / whitelist
 - object / column
 - generates list of sqls
 - replications to other db platforms possible

Replication



- Self Repair / Healing after Outtages
 - Iog stream to extract / apply gaps
- Management by GUI
- CDC / Streams alternative
- Parallel Capture and Apply
- EE or Partitioning not necessary

Presentation References



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